

Tutorial 2 - Solutions

Exercise 1*

(Assume that A, B are process constants and a, b are channel names.)

- $a.b.A + B$ **Correct**
- $(a.Nil + \bar{a}.A) \setminus \{a, b\}$ **Correct**
- $(a.Nil | \bar{a}.A) \setminus \{a, \tau\}$ **False**, τ can not be used in a restriction
- $a.B + [a/b]$ **False**, relabelling can be applied only on a valid process expression
- $\tau.\tau.B + Nil$ **Correct**
- $(a.B + b.B)[a/b, b/a]$ **Correct**
- $(a.B + \tau.B)[a/\tau, b/a]$ **False**, the relabeling function should satisfy $f(\tau) = \tau$ but here $f(\tau) = a$
- $(a.B + \tau.B)[\tau/a]$ **Correct**, any action can be relabelled to τ
- $(a.b.A + \bar{a}.Nil) | B$ **Correct**
- $(a.b.A + \bar{a}.Nil).B$ **False**, only actions can be used as prefixes
- $(a.b.A + \bar{a}.Nil) + B$ **Correct**
- $(Nil | Nil) + Nil$ **Correct**

Exercise 2*

- Derivation of $(A | \bar{b}.Nil) \setminus \{b\} \xrightarrow{\tau} (a.B | Nil) \setminus \{b\}$.

$$\begin{array}{c}
 \text{ACT} \frac{}{b.a.B \xrightarrow{b} a.B} \\
 \text{CON} \frac{}{A \xrightarrow{b} a.B} \\
 \text{COM3} \frac{}{(A | \bar{b}.Nil) \xrightarrow{\tau} (a.B | Nil)} \\
 \text{RES} \frac{}{(A | \bar{b}.Nil) \setminus \{b\} \xrightarrow{\tau} (a.B | Nil) \setminus \{b\}} \quad \tau, \bar{\tau} \notin \{b\}
 \end{array}$$

- Derivation of $(A | \bar{b}.a.B) + (\bar{b}.A)[a/b] \xrightarrow{\bar{b}} (A | a.B)$.

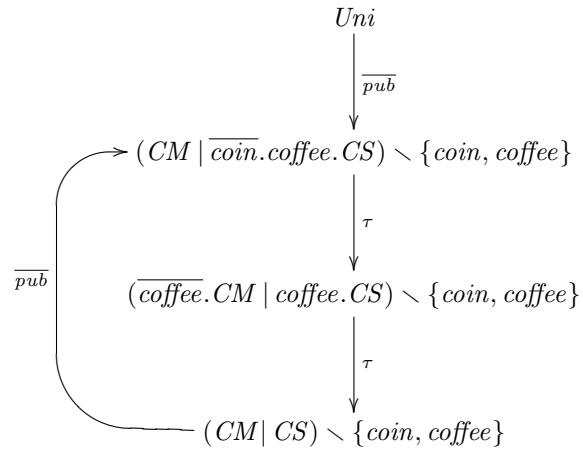
$$\begin{array}{c}
 \text{ACT} \frac{}{\bar{b}.a.B \xrightarrow{\bar{b}} a.B} \\
 \text{COM2} \frac{}{A | \bar{b}.a.B \xrightarrow{\bar{b}} A | a.B} \\
 \text{SUM1} \frac{}{(A | \bar{b}.a.B) + (\bar{b}.A)[a/b] \xrightarrow{\bar{b}} (A | a.B)}
 \end{array}$$

- Derivation of $(A | \bar{b}.a.B) + (\bar{b}.A)[a/b] \xrightarrow{\bar{a}} A[a/b]$.

$$\begin{array}{c}
 \text{ACT} \frac{}{\bar{b}.A \xrightarrow{\bar{b}} A} \\
 \text{REL} \frac{}{(\bar{b}.A)[a/b] \xrightarrow{\bar{a}} A[a/b]} \\
 \text{SUM2} \frac{}{(A | \bar{b}.a.B) + (\bar{b}.A)[a/b] \xrightarrow{\bar{a}} A[a/b]}
 \end{array}$$

Exercise 3*

LTS for the process $Uni \stackrel{\text{def}}{=} (CM \mid CS) \setminus \{coin, coffee\}$.

**Exercise 4**

Transition system for $A \stackrel{\text{def}}{=} (a.A) \setminus \{b\}$.

$$A \xrightarrow{a} A \setminus \{b\} \xrightarrow{a} (A \setminus \{b\}) \setminus \{b\} \xrightarrow{a} ((A \setminus \{b\}) \setminus \{b\}) \setminus \{b\} \xrightarrow{a} \dots$$

One solution could be the CCS defining equation $B \stackrel{\text{def}}{=} a.B$ which generates a finite LTS with (intuitively) the same behavior as A .