

# Final coalgebras in categories with factorization systems

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**Abstract.** For a functor  $T: \mathcal{C} \rightarrow \mathcal{C}$ , we show that, if  $\mathcal{C}$  admits a factorization system  $(\mathcal{L}, \mathcal{R})$  with all arrows in  $\mathcal{R}$  monic and  $\mathcal{C}$  is  $\mathcal{R}$ -well-powered, then the final coalgebra is characterized constructively as the  $\mathcal{R}$ -union of certain sets of  $T$ -coalgebras, provided that the final sequence of  $T$  has an  $\mathcal{R}$ -arrow at some ordinal  $\alpha$ , and  $T$  preserves  $\mathcal{R}$ -morphisms.

**Final sequence:** Let  $T$  be an endofunctor on a category  $\mathcal{C}$  with final object and limits of ordinal-indexed diagrams. The *final sequence* of  $T$  is a limit-preserving functor  $A: \mathbf{Ord}^{op} \rightarrow \mathcal{C}$  such that, for all ordinals  $\gamma \leq \beta$ ,  $A(\beta+1) = TA(\beta)$ ,  $A(\beta+1 \rightarrow \gamma+1) = TA(\beta \rightarrow \gamma)$ , and  $A(0) = 1$ . Note that, since  $A$  preserves limits, for all limit ordinals  $\beta$ , the arrow  $A(\beta) \rightarrow \lim_{\gamma < \beta} A(\gamma)$  is an isomorphism.

In [1,2], it is shown that if this sequence *stabilizes* at some  $\alpha$ , in the sense that  $f = A(\alpha+1 \rightarrow \alpha)$  is an isomorphism, then  $(A(\alpha), f^{-1})$  is a final  $T$ -coalgebra. This follows since, for any  $T$ -coalgebra  $(X, h)$  and ordinal  $\beta$ , there exists a cone  $(X, (h_\gamma)_{\gamma \in \beta^{op}})$  on  $A \upharpoonright \beta^1$  uniquely determined by  $A(\gamma+1 \rightarrow \gamma) \circ Th_\gamma \circ h = h_\gamma$ , for all  $\gamma \leq \beta$ . Therefore,  $h_\alpha: (X, h) \rightarrow (A(\alpha), f^{-1})$  is a morphism of  $T$ -coalgebras, which is easily seen to be unique. Note that, this method is constructive if one can determine an ordinal  $\alpha$  at which the final sequence stabilizes.

In [3], Worrell shows that for any mono-preserving accessible endofunctor on a locally presentable category the final sequence stabilizes. However, the proof does not give any constructive bound for stabilization. If one restricts the attention to only  $\kappa$ -accessible endofunctors on **Set**,  $\kappa + \kappa$  steps are sufficient for the final sequence to stabilize. This bound depends heavily on the fact that in **Set**, all monomorphisms split, which is indeed a very strong requirement.

Both the results need that the final sequence  $A$  reaches a monic arrow at some  $\alpha$ , then stabilization follows since the category is well-powered and all  $A(\gamma)$  are subobjects of  $A(\alpha)$ , for all  $\gamma \geq \alpha$ . The restriction on accessible functors on locally accessible categories ensures that these requirements hold, and that the underlying category has a (strong-epi, mono) factorization system.

**A new characterization:** Let  $(\mathcal{L}, \mathcal{R})$  be a factorization system for  $\mathcal{C}$ , such that arrows in  $\mathcal{R}$  are monic, and let  $\mathcal{C}$  be  $\mathcal{R}$ -well-powered. Under these hypotheses, we give a characterization of a final coalgebra for a functor  $T: \mathcal{C} \rightarrow \mathcal{C}$  via its final sequence. The use of the final sequence is twofold: it guarantees unicity of

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<sup>1</sup>  $A \upharpoonright \beta: \beta^{op} \rightarrow \mathcal{C}$  restricts  $A$  on the full subcategory of  $\mathbf{Ord}^{op}$  of all ordinals  $\gamma \leq \beta$ .

the final homomorphism, and provides a weakly final coalgebra. Notably, we do not need any bound for stabilization, still the proof is constructive.

**Theorem 1.** *Assume  $A$ , the final sequence of  $T$ , is such that  $A(\alpha+1 \rightarrow \alpha) \in \mathcal{R}$ , for some  $\alpha$ , and that  $T$  preserves  $\mathcal{R}$ -morphisms. Then, for any  $T$ -coalgebra  $(X, h)$  and  $\rho_h \circ \lambda_h$   $(\mathcal{L}, \mathcal{R})$ -factorization of  $h_\alpha$ , there exists a unique arrow  $\phi_h$  making the diagram aside commute.*

$$\begin{array}{ccccc} X & \xrightarrow{\lambda_h} \triangleright & F_h & \xrightarrow{\rho_h} & A(\alpha) \\ h \downarrow & & \downarrow \phi_h & & \uparrow \\ TX & \xrightarrow{T\lambda_h} & TF_h & \xrightarrow{T\rho_h} & A(\alpha+1) \end{array}$$

Moreover, if  $(X, h)$  is weakly final, then  $(F_h, \phi_h)$  is a final  $T$ -coalgebra.

*Proof.* (Sketch) Since  $T$  preserves  $\mathcal{R}$ -morphisms,  $T\rho_h \in \mathcal{R}$ . Therefore, the outer square diagram is a lifting problem for  $\lambda_h \in \mathcal{L}$  and  $A(\alpha+1 \rightarrow \alpha) \circ T\rho_h \in \mathcal{R}$ , and  $\phi_h$  is solution to it. Assume  $(X, h)$  is weakly final, then  $(F_h, \phi_h)$  is weakly final too, since  $\lambda_h$  is a  $T$ -homomorphism. Unicity follows by left cancellability of  $\rho_h$ , since, for any  $T$ -coalgebra  $(Y, k)$  and arrow  $f: Y \rightarrow A(\alpha)$  such that  $f = A(\alpha+1 \rightarrow \alpha) \circ Tf \circ k$ , one proves by transfinite induction that  $f = k_\alpha$ .  $\square$

Note that,  $(\mathcal{L}, \mathcal{R})$ -factorizations of morphisms are not unique, hence for a given  $T$ -coalgebra  $(X, h)$ , the associated  $T$ -coalgebra  $(F_h, \phi_h)$  is not uniquely determined. However, under the hypothesis of Theorem 1, one can fix any factorization  $h_\alpha = \rho_h \circ \lambda_h$  to obtain an endofunctor  $F$  on the category of  $T$ -coalgebras, mapping objects  $(X, h)$  to  $(F_h, \phi_h)$ , and morphisms  $f: (X, h) \rightarrow (Y, k)$  to the unique solution  $\varphi_f$  of the lifting problem on the right. Functoriality crucially depends on the assumption that all  $\mathcal{R}$ -morphisms are monic (see the Appendix for details).

$$\begin{array}{ccccc} X & \xrightarrow{\lambda_h} \triangleright & F_h & \xrightarrow{\rho_h} & A(\alpha) \\ f \downarrow & & \downarrow \varphi_f & & \parallel \\ Y & \xrightarrow{\lambda_k} \triangleright & F_k & \xrightarrow{\rho_k} & A(\alpha) \end{array}$$

Finally, observe that, for all  $T$ -coalgebras  $(X, h)$ ,  $F_h$  is an  $\mathcal{R}$ -subobject of  $A(\alpha)$ , and since  $\mathcal{C}$  is assumed to be  $\mathcal{R}$ -well-powered, there must be only a set  $I$  (up to isomorphism) of such  $F_h$ 's. Thus, if  $\mathcal{C}$  has coproducts, we are allowed to take the coproduct coalgebra  $\coprod_I (F_i, \phi_i)$ , which is readily seen to be weakly final, with homomorphism from any  $T$ -coalgebra  $(X, h)$  given by

$$X \xrightarrow{\lambda_h} F_h \xrightarrow{\cong} F_i \xrightarrow{in_i} \coprod F_i$$

where  $F_i$  is the representative of  $F_h$  in  $I$ .

Applying Theorem 1 to  $\coprod_I (F_i, \phi_i)$ , the final  $T$ -coalgebra is just the  $\mathcal{R}$ -union of the coalgebras in  $I$ . Notably, finality does not depend on the choice of  $I$ , which can be determined constructively by an analysis on the  $\mathcal{R}$ -subobjects of  $A(\alpha)$ .

## References

1. J. Adámek and V. Koubek. On the greatest fixed point of a set functor. *Theoretical Comput. Sci.*, 150(1):57–75, 1995.
2. Michael and Barr. Algebraically compact functors. *Journal of Pure and Applied Algebra*, 82(3):211 – 231, 1992.
3. J. Worrell. Terminal sequences for accessible endofunctors. *Electronic Notes in Theoretical Computer Science*, 19:24–38, 1999.

## Appendix

Here we provide full proofs of all the technical statements mentioned above.

**Lemma 2.** *Let  $A$  be the final sequence of  $T$ . For any  $T$ -coalgebra  $(X, h)$  and ordinal  $\alpha$ , if a morphism  $k: X \rightarrow A(\alpha)$  is such that  $k = A(\alpha+1 \rightarrow \alpha) \circ Tk \circ h$ , then  $k = h_\alpha$ .*

*Proof.* For  $\beta \leq \alpha$ , let  $k_\beta = A(\alpha \rightarrow \beta) \circ k$ . We will show by transfinite induction that  $k_\beta = h_\beta$ , for  $\beta \leq \alpha$ . Certainly  $k_0 = h_0$  since their codomain is the final object. Assuming  $k_\beta = h_\beta$ , we have

$$\begin{aligned}
 k_{\beta+1} &= A(\alpha \rightarrow \beta+1) \circ k && \text{(by def. } k_{\beta+1}\text{)} \\
 &= A(\alpha \rightarrow \beta+1) \circ A(\alpha+1 \rightarrow \alpha) \circ Tk \circ h && \text{(by hp.)} \\
 &= A(\alpha+1 \rightarrow \beta+1) \circ Tk \circ h && \text{(by funct. } A\text{)} \\
 &= TA(\alpha \rightarrow \beta) \circ Tk \circ h && \text{(by def. } A\text{)} \\
 &= Tk_\beta \circ h && \text{(by def. of } k_\beta\text{)} \\
 &= Th_\beta \circ h && \text{(by inductive hp.)} \\
 &= h_{\beta+1} && \text{(by def. } h_{\beta+1}\text{)}
 \end{aligned}$$

Suppose  $\beta$  is a limit ordinal and  $k_\gamma = h_\gamma$ , for every  $\gamma < \beta$ . By definition of  $k_\gamma$  and by the fact that  $(X, (h_\gamma)_{\gamma \in \alpha^{op}})$  is a cone on  $A \upharpoonright \alpha$ ,  $A(\alpha \rightarrow \gamma) \circ k = A(\alpha \rightarrow \gamma) \circ h_\alpha$ . Since  $A(\beta) \cong \lim_{\gamma < \beta} A(\gamma)$ , we have  $A(\alpha \rightarrow \beta) \circ k = A(\alpha \rightarrow \beta) \circ h_\alpha$ . Therefore, by definition of  $k_\beta$  and compatibility of the cone,  $k_\beta = h_\beta$ .  $\square$

*Proof.* (Theorem 1) By hypothesis  $T$  preserves  $\mathcal{R}$ -morphisms, hence  $T\rho_h \in \mathcal{R}$ . Therefore  $A(\alpha+1 \rightarrow \alpha) \circ T\rho_h \in \mathcal{R}$ , since by hypothesis  $A(\alpha+1 \rightarrow \alpha) \in \mathcal{R}$  and  $\mathcal{R}$  is closed by composition. By definition,  $h_\alpha = A(\alpha+1 \rightarrow \alpha) \circ Th_\alpha \circ h$ , therefore, the diagram below is a lifting problem for  $\lambda_h \in \mathcal{L}$  and  $A(\alpha+1 \rightarrow \alpha) \circ T\rho_h \in \mathcal{R}$

$$\begin{array}{ccccc}
 & & \xrightarrow{h_\alpha} & & \\
 & & \curvearrowright & & \\
 X & \xrightarrow{\lambda_h} & F_h & \xrightarrow{\rho_h} & A(\alpha) \\
 & & \downarrow \phi_h & & \uparrow \\
 h \downarrow & & & & \\
 TX & \xrightarrow{T\lambda_h} & TF_h & \xrightarrow{T\rho_h} & A(\alpha+1) \\
 & & \curvearrowleft & & \\
 & & \xrightarrow{Th_\alpha} & & 
 \end{array}$$

hence,  $\phi_h$  can be chosen as a solution to it, so that all the sub-diagrams commute.

Assume  $(X, h)$  is weakly final. We will prove that  $(F_h, \phi_h)$  is a final  $T$ -coalgebra. Let  $(Y, k)$  be a  $T$ -coalgebra, then, by weakly finality of  $(X, h)$ , there exists a morphism  $f: (Y, k) \rightarrow (X, h)$ . Note that,  $\lambda_h: (X, h) \rightarrow (F_h, \phi_h)$  is a morphism of  $T$ -coalgebras, therefore  $\lambda_h \circ f$  is a  $T$ -homomorphism from  $(Y, k)$  to  $(F_h, \phi_h)$ . This proves weak finality. As for unicity, let  $f, g: (Y, h) \rightarrow (F_h, \phi_h)$  be two morphisms of  $T$ -coalgebras. Consider the two composites  $\lambda_h \circ f$  and  $\lambda_h \circ g$ . Since both composites satisfy the conditions of Lemma 2, we have that

$\lambda_h \circ f = \lambda_h \circ g$ . Now, since  $\lambda_h \in \mathcal{R}$  and all  $\mathcal{R}$ -morphisms are monic, by left cancellability of monomorphisms we conclude that  $f = g$ .  $\square$

Another interesting consequence of the fact that all arrows in  $\mathcal{R}$  are monic is that the mapping  $(X, h) \mapsto (F_h, \phi_h)$  in Theorem 1, can be extended to a functor.

**Proposition 3.** *Let  $A$  be the final sequence of  $T$ . Assume that  $T$  preserves  $\mathcal{R}$ -morphisms and that  $A(\alpha+1 \rightarrow \alpha) \in \mathcal{R}$ , for some  $\alpha$ . Then, we can define a functor  $F: T\text{-Coalg} \rightarrow T\text{-Coalg}$  as follows:*

**Objects:** *For a  $T$ -coalgebra  $(X, h)$ , we define  $F(X, h) = (F_h, \phi_h)$ , where  $\rho_h \circ \lambda_h$  is a distinguished  $(\mathcal{L}, \mathcal{R})$ -factorization of  $h_\alpha$ , and  $\phi_h$  is the (unique) solution to the lifting problem below.*

$$\begin{array}{ccccc} X & \xrightarrow{\lambda_h} \triangleright & F_h & \xrightarrow{\rho_h} & A(\alpha) \\ h \downarrow & & \downarrow \phi_h & & \uparrow \\ TX & \xrightarrow{T\lambda_h} & TF_h & \xrightarrow{T\rho_h} & A(\alpha+1) \end{array}$$

**Morphisms** *For a morphism  $f: (X, h) \rightarrow (Y, k)$  of  $T$ -coalgebras, we define  $Ff = \varphi_f$ , where  $\varphi_f$  is the (unique) solution to the lifting problem below.*

$$\begin{array}{ccccc} X & \xrightarrow{\lambda_h} \triangleright & F_h & \xrightarrow{\rho_h} & A(\alpha) \\ f \downarrow & & \downarrow \varphi_f & & \parallel \\ Y & \xrightarrow{\lambda_k} \triangleright & F_k & \xrightarrow{\rho_k} & A(\alpha) \end{array}$$

*Proof.* We first check that  $F$  is well defined. Let  $f: (X, h) \rightarrow (Y, k)$  be a morphism of  $T$ -coalgebras, then we have to prove that  $Ff = \varphi_f: F_h \rightarrow F_k$  is a homomorphism of  $T$ -coalgebras. By definition, we have

$$\begin{aligned} A(\alpha+1 \rightarrow \alpha) \circ T\rho_k \circ T\varphi_f \circ \phi_h &= A(\alpha+1 \rightarrow \alpha) \circ T\rho_h \circ \phi_h && \text{(by def. } T\varphi_f) \\ &= \rho_h && \text{(by def. } \phi_h) \\ &= \rho_k \circ \varphi_f && \text{(by def. } \varphi_f) \\ &= A(\alpha+1 \rightarrow \alpha) \circ T\rho_k \circ \phi_k \circ \varphi_f && \text{(by def. } \rho_k) \end{aligned}$$

Since  $T$  preserves  $\mathcal{R}$ -morphisms and  $\mathcal{R}$  is closed under composition, we have that  $A(\alpha+1 \rightarrow \alpha) \circ T\rho_k \in \mathcal{R}$ , therefore is monic. By left cancellability of monomorphisms, we have  $T\varphi_f \circ \phi_h = \phi_k \circ \varphi_f$ , hence  $\varphi_f$  is a  $T$ -homomorphism.

Factoriality is easily proved. Let  $f: (X, h) \rightarrow (Y, k)$  and  $g: (Y, k) \rightarrow (Z, l)$  be morphisms in  $T\text{-Coalg}$ . By definition  $\rho_h = \rho_l \circ F(g \circ f)$ ,  $\rho_h = \rho_k \circ F(f)$ , and  $\rho_k = \rho_l \circ F(g)$ , therefore  $\rho_l \circ F(g \circ f) = \rho_l \circ F(g) \circ F(f)$ . By left cancellability of  $\rho_l$  we have  $F(g \circ f) = F(g) \circ F(f)$ . Consider  $Fid_{(X, h)}$ . By definition, we have  $\rho_h \circ Fid_{(X, h)} = \rho_h$ , therefore  $\rho_h \circ Fid_{(X, h)} = \rho_h \circ id_X$ . Again, by left cancellability of  $\rho_h$ , we have  $Fid_{(X, h)} = id_X$ , and since  $id_X = id_{(X, h)}$ , we have  $Fid_{(X, h)} = id_{(X, h)}$ .  $\square$