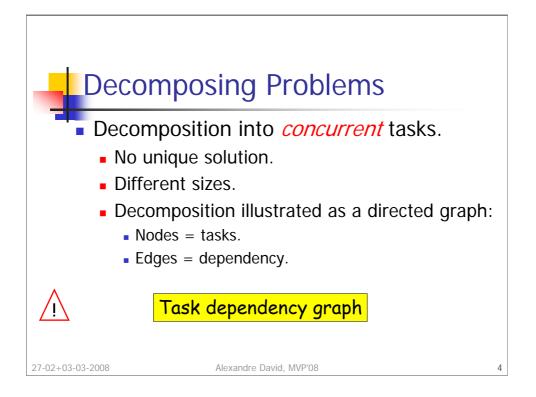


There are other courses specifically on concurrency. We won't treat the problems proper to concurrency such as deadlocks, livelocks, theory on semaphores and synchronization. However, we will use them, and when needed, apply techniques to avoid problems like deadlocks.

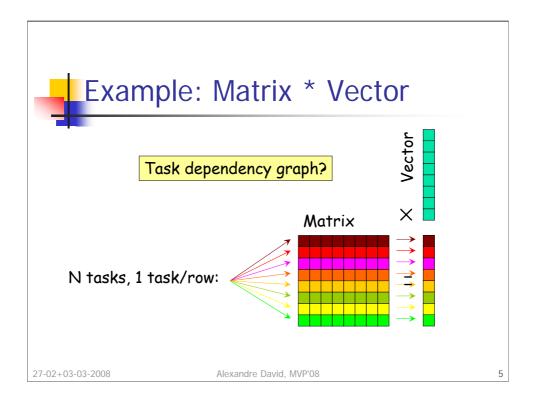


Many solutions are often possible but few will yield good performance and be scalable. We have to consider the computational and storage resources needed to solve the problems.

Size of the tasks in the sense of the amount of work to do. Can be more, less, or unknown. Unknown in the case of a search algorithm is common.

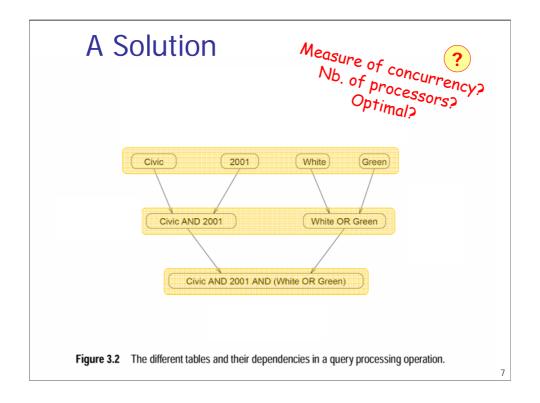
Dependency: All the results from incoming edges are required for the tasks at the current node.

We will not consider tools for automatic decomposition. They work fairly well only for highly structured programs or options of programs.

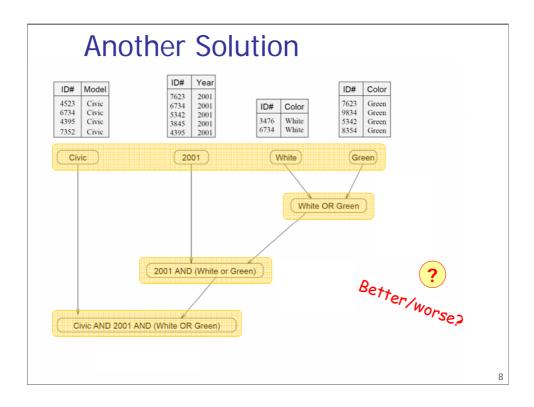


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(CO	LOR = `	`GRE	EN''	OR CO	LOR = `	`WHITE)	
ID#	Model	Year	Color	Dealer	Price		
4523	Civic	2002	Blue	MN	\$18,000		
3476	Corolla	1999	White	IL	\$15,000		
7623	Camry	2001	Green	NY	\$21,000		
9834	Prius	2001	Green	CA	\$18,000		
6734	Civic	2001	White	OR	\$17,000		
5342	Altima	2001	Green	FL	\$19,000	-	
3845	Maxima	2001	Blue	NY	\$22,000		
8354	Accord	2000	Green	VT	\$18,000		
4395	Civic	2001	Red	CA	\$17,000		
7352	Civic	2002	Red	WA	\$18,000		

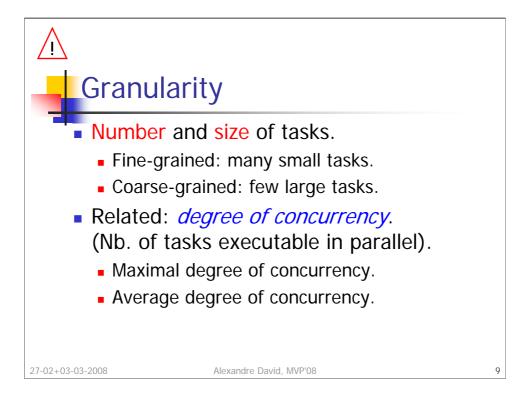
The question is: How to decompose this into concurrent tasks? Different tasks may generate intermediate results that will be used by other tasks.



How much concurrency do we have here? How many processors to use? Is it optimal?



Is it better or worse? Why?



•Previous matrix*vector fine-grained.

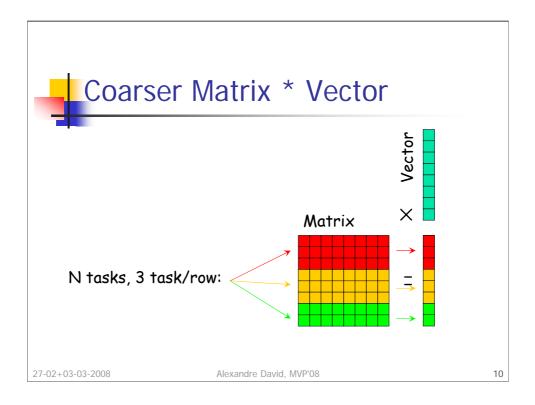
•Database example coarse grained.

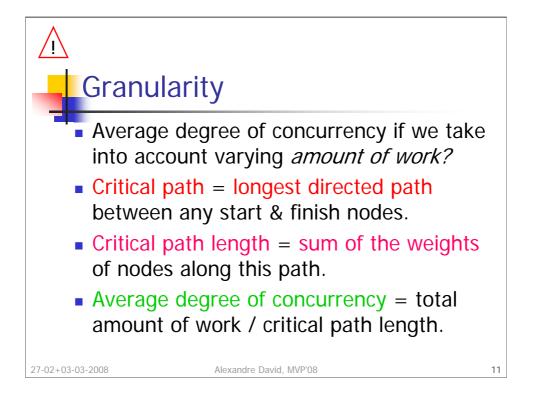
Degree of concurrency: Number of tasks that can be executed in parallel.

Average degree of concurrency is a more useful measure.

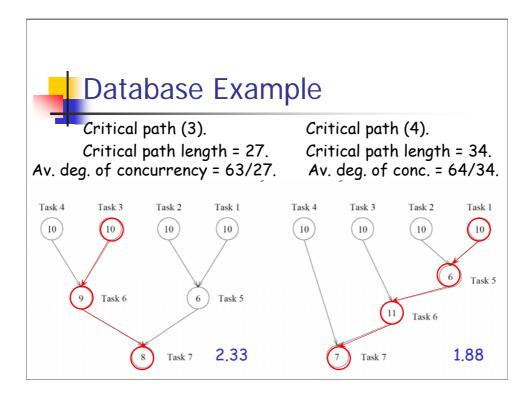
Assume that the tasks in the previous database examples have the same granularity. What's their average degrees of concurrency? 7/3=2.33 and 7/4=1.75.

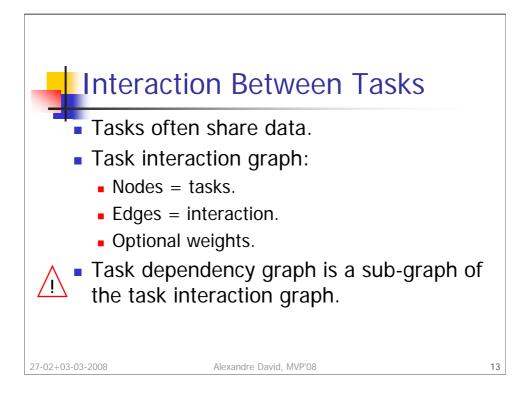
Common sense: Increasing the granularity of decomposition and utilizing the resulting concurrency to perform more tasks in parallel increases performance. However, there is a limit to granularity due to the nature of the problem itself.





Weights on nodes denote the amount of work to be done on these nodes. Longest path \rightarrow shortest time needed to execute in parallel.



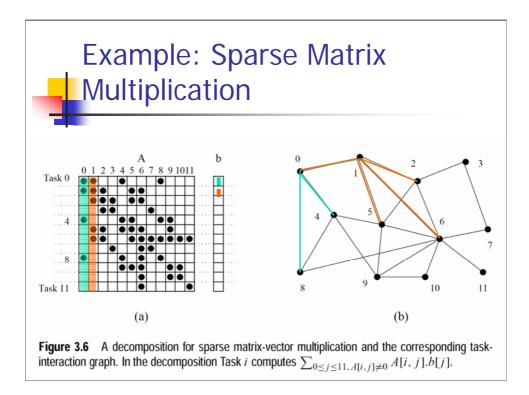


Another important factor is interaction between tasks on different processors.

Share data implies synchronization protocols (mutual exclusion, etc) to ensure **consistency**.

Edges generally undirected. When directed edges are used, they show the direction of the flow of data (and the flow is unidirectional).

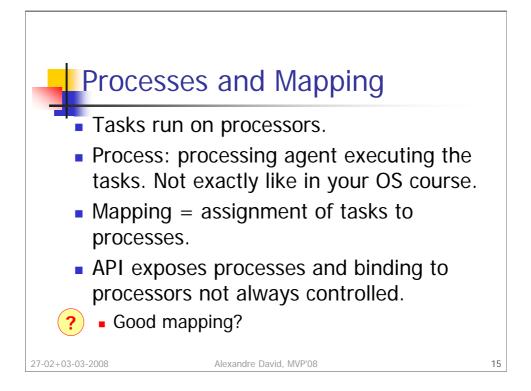
Dependency between tasks implies interaction between them.



Sparse matrix: A significant number of its entries are zero and the zeros do not conform to predefined patterns. Typically, we do not need to take the zeros into account.

In the example: Task i owns row i of A and b.

Interaction depends on the mapping work to do / task, i.e., granularity, and mapping tasks – processor.



Here we are not talking directly on the mapping to processors. A processor can execute two processes.

Good mapping:

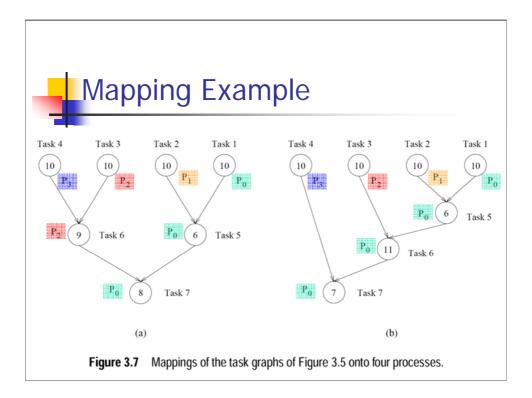
•Maximize concurrency by mapping independent tasks to different processes.

•Minimize interaction by mapping interacting tasks on the same process.

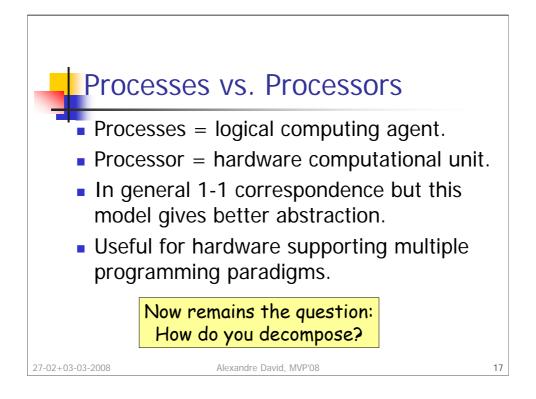
Can be conflicting, good trade-off is the key to performance.

Decomposition determines degree of concurrency.

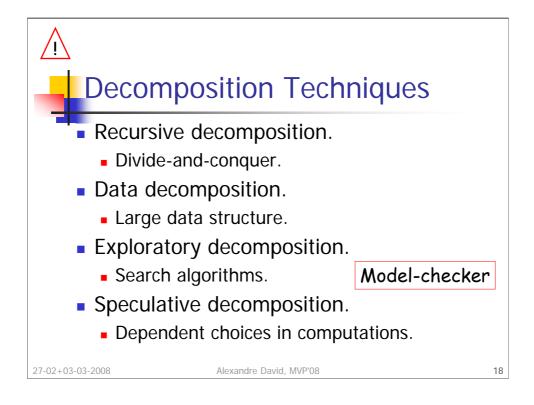
Mapping determines how much concurrency is utilized and how efficiently.

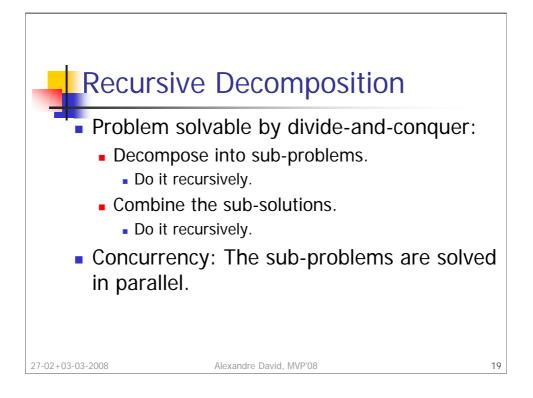


Notice that the mapping keeps one process from the previous stage because of dependency: We can avoid interaction by keeping the same process.



- Example of hybrid hardware: cluster of MP machines. Each node has shared memory and communicates with other nodes via MPI.
- 1. Decompose and map to processes for MPI.
- 2. Decompose again but suitable for shared memory.



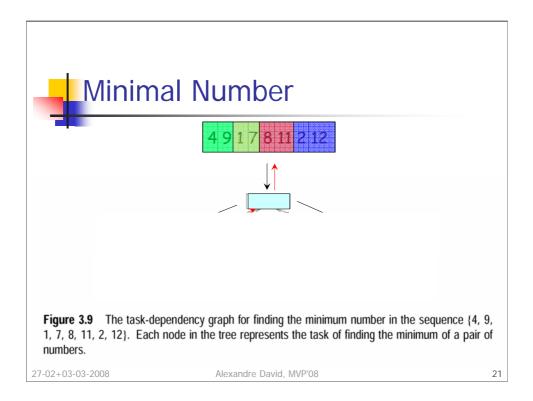


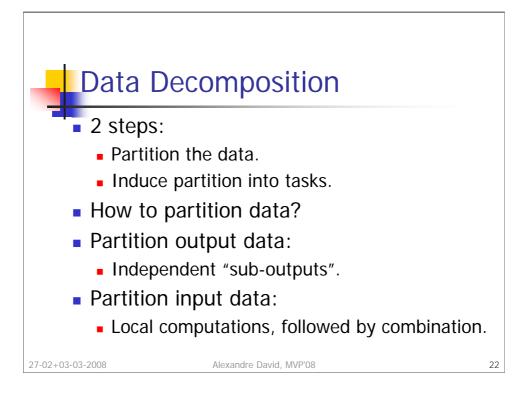
Small problem is to start and finish: with one process only.

Quicksort Example	
5 12 11 1 10 6 8 3 7 4 9 2	
Figure 3.8 The quicksort task-dependency graph based on recursive decomposition for sorting a sequence of 12 numbers.	1

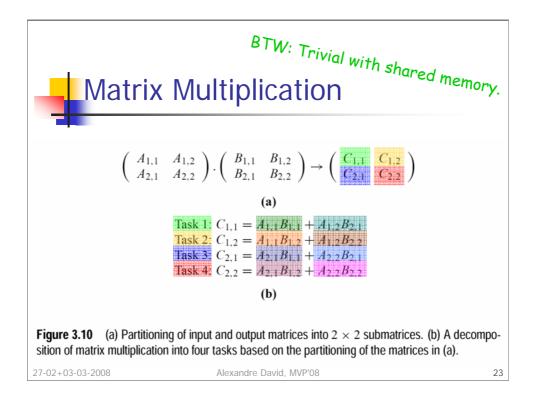
Recall on the quicksort algorithm:

- •Choose a pivot.
- •Partition the array.
- •Recursive call.
- •Combine result: nothing to do.

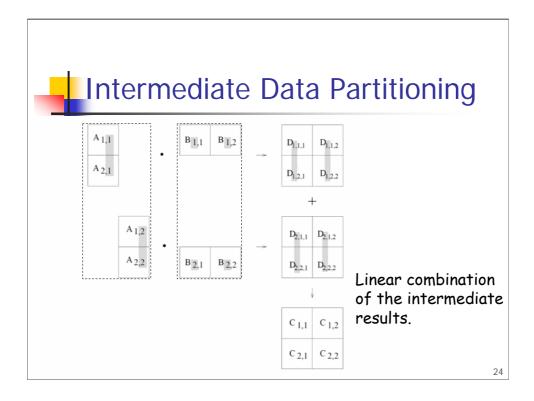


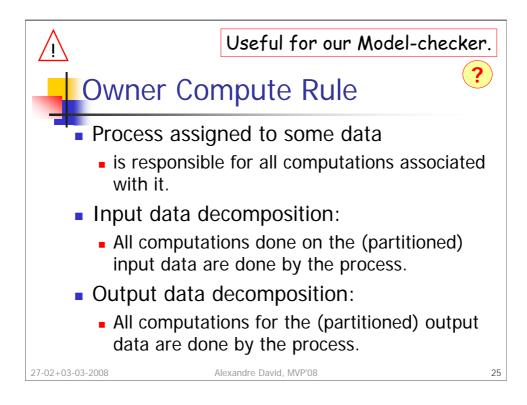


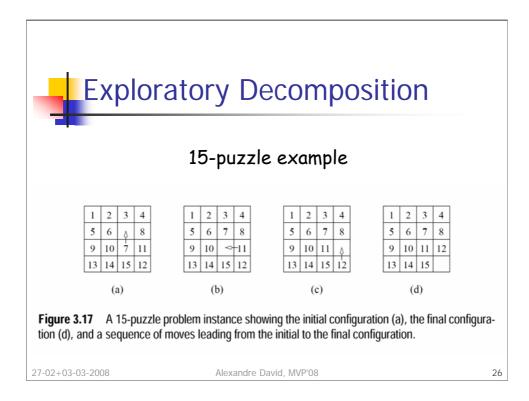
Partitioning of input data is a bit similar to divide-and-conquer.



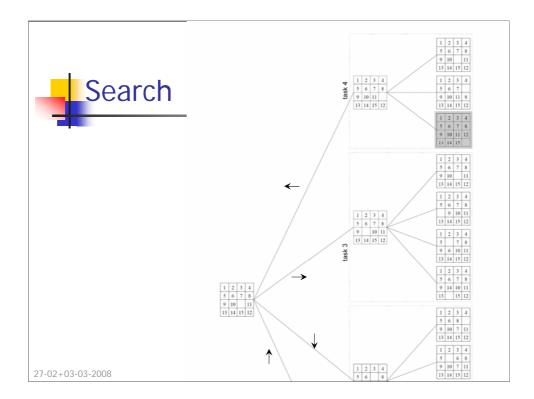
We can partition further for the tasks. Notice the dependency between tasks. What is the task dependency graph?

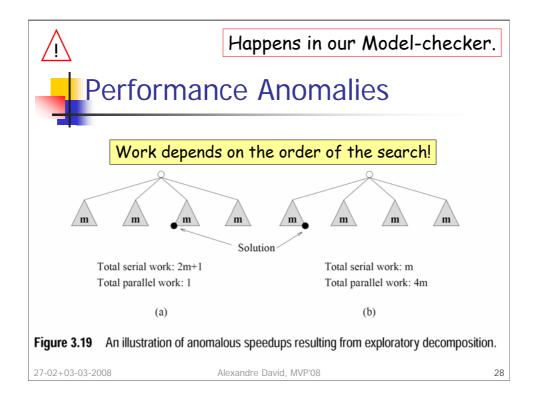


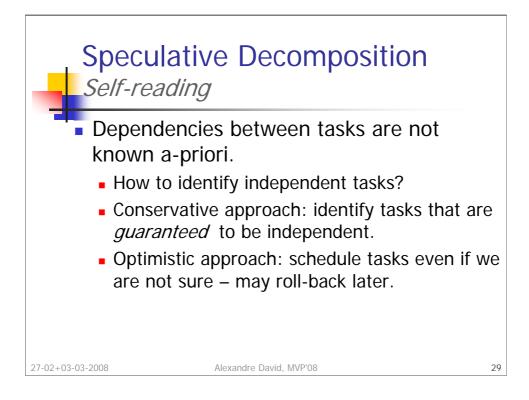




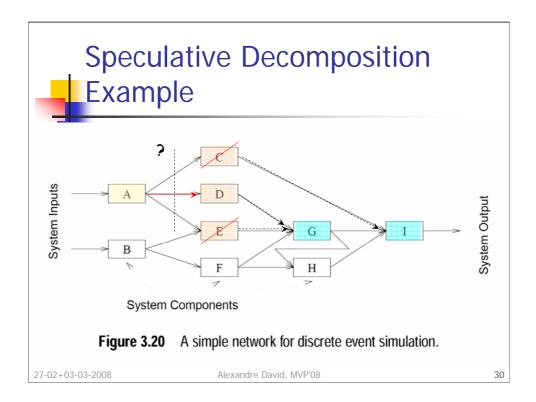
Suitable for search algorithms. Partition the search space into smaller parts and search in parallel. We search the solution by a tree search technique.







Not possible to identify independent tasks in advance. Conservative approaches may yield limited concurrency. Optimistic approach = speculative. Optimistic approach is similar to branch prediction algorithms in processors.



More aggregate work is done. Problem is to send inputs to the next stages speculatively. Could be the case that two different kinds of outputs are possible for A and A could start C,D,E twice.

Other approaches are possible that combine different techniques: hybrid decompositions.