Sorting (Chapter 9)

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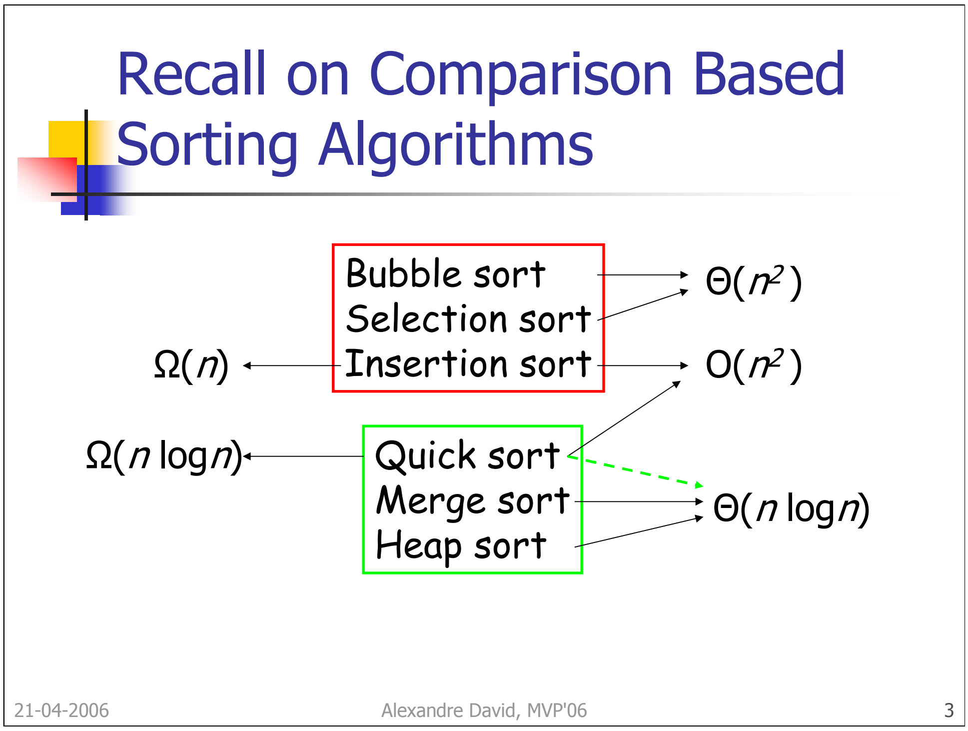
B2-206

Sorting

Problem

Arrange an unordered collection of elements into monotonically increasing (or decreasing) order.
Let $S = \langle a_1, a_2, \dots, a_n \rangle$.
Sort S into $S' = \langle a_1', a_2', \dots, a_n' \rangle$ such that
$a_i' \leq a_j'$ for $1 \leq i \leq j \leq n$
and S' is a permutation of S.

The elements to sort (actually used for comparisons) are also called the keys.



You should know these complexities from a previous course on algorithms.

Characteristics of Sorting Algorithms

- In-place sorting: No need for additional memory (or only constant size).
- Stable sorting: Ordered elements keep their original relative position.
- Internal sorting: Elements fit in process memory.
- External sorting: Elements are on auxiliary storage.

We assume internal sorting is possible.

Fundamental Distinction

- **Comparison based** sorting:
  - *Compare-exchange* of pairs of elements.
  - Lower bound is $\mathbf{\Omega(n \log n)}$ (proof based on decision trees).
  - Merge & heap-sort are optimal.
- **Non-comparison based** sorting:
  - Use information on the element to sort.
  - Lower bound is $\mathbf{\Omega(n)}$.
  - Counting & radix-sort are optimal.

We assume comparison based sorting is used.

Issues in Parallel Sorting

- Where to store input & output?
  - One process or distributed?
  - Enumeration of processes used to distribute output.
- How to compare?
  - How many elements per process?
  - As many processes as element $\Rightarrow$ poor performance because of inter-process communication.

Parallel Compare-Exchange

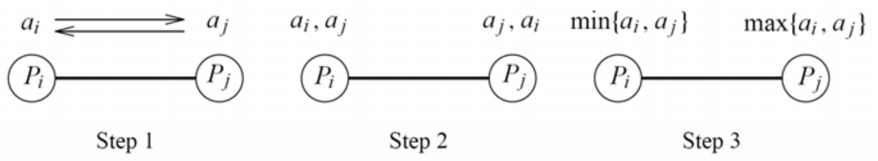
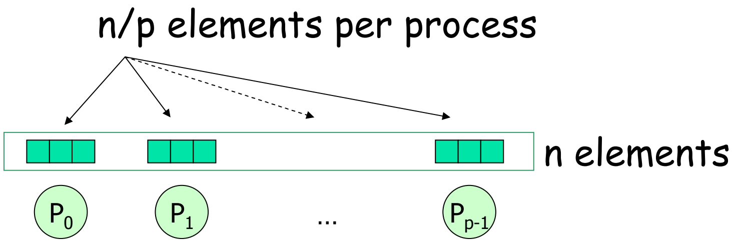


Figure 9.1 A parallel compare-exchange operation. Processes P_i and P_j send their elements to each other. Process P_i keeps $\min\{a_i, a_j\}$, and P_j keeps $\max\{a_i, a_j\}$.

Communication cost: t_s+t_w.
Comparison cost much cheaper $\Rightarrow$ communication time dominates.

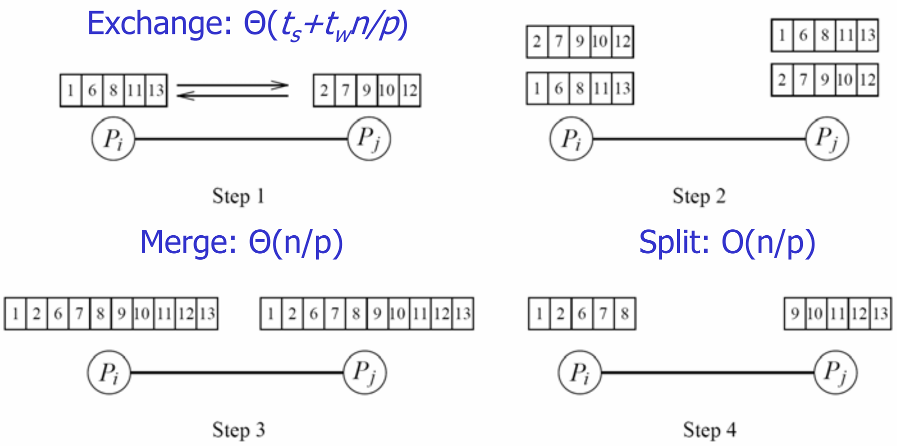
Blocks of Elements Per Process



Blocks: $A_0 \leq A_1 \leq \ldots \leq A_{p-1}$

Compare-Split

For large blocks: $\Theta(n/p)$



Sorting Networks

- Mostly of theoretical interest.
- Key idea: Perform many comparisons in parallel.
- Key elements:
  - Comparators: 2 inputs, 2 outputs.
  - Network architecture: Comparators arranged in columns, each performing a permutation.
  - Speed proportional to the depth.

Comparators

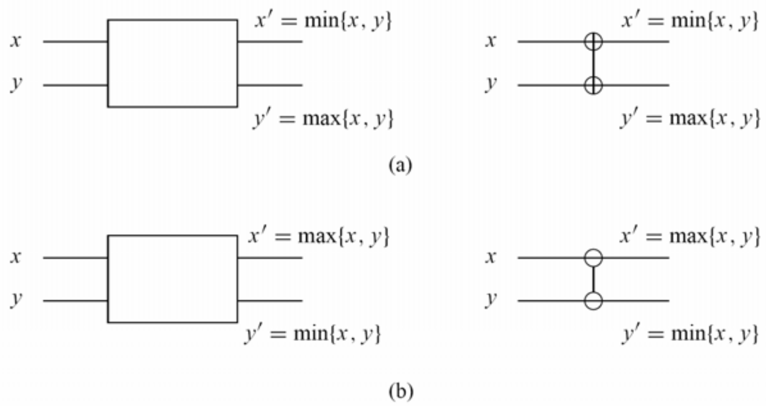


Figure 9.3 A schematic representation of comparators: (a) an increasing comparator, and (b) a decreasing comparator.

Sorting Networks

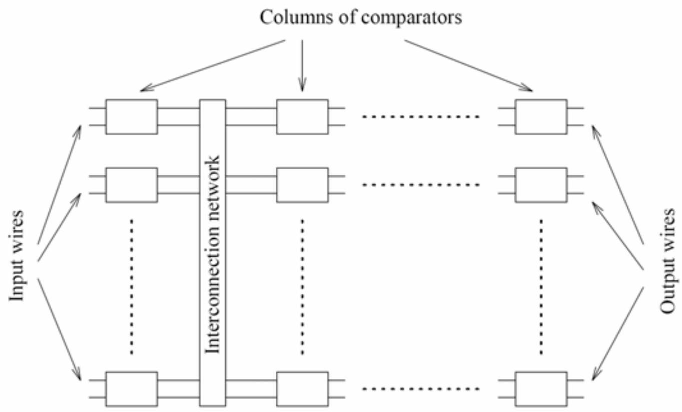


Figure 9.4 A typical sorting network. Every sorting network is made up of a series of columns, and each column contains a number of comparators connected in parallel.

Bitonic Sequence

Definition

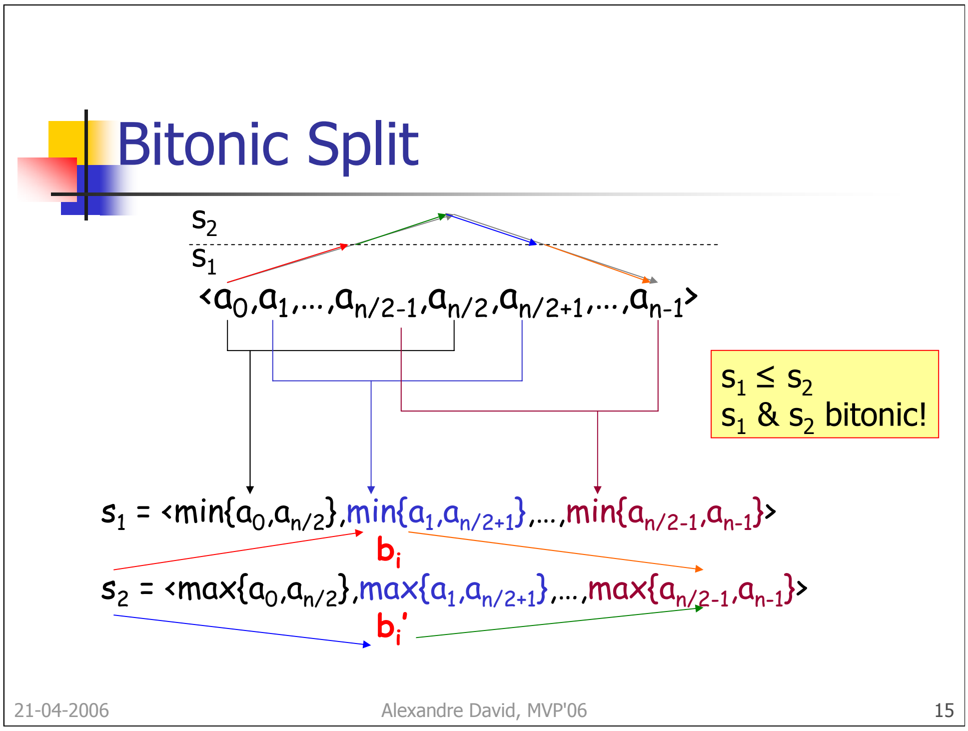
A **bitonic sequence** is a sequence of elements $\langle a_0,a_1,\ldots,a_n\rangle$ s.t.

1. $\exists i$, $0 \le i \le n-1$ s.t. $\langle a_0,\ldots,a_i\rangle$ is monotonically increasing and $\langle a_{i+1},\ldots,a_{n-1}\rangle$ is monotonically decreasing,
2. or there is a cyclic shift of indices so that 1) is satisfied.

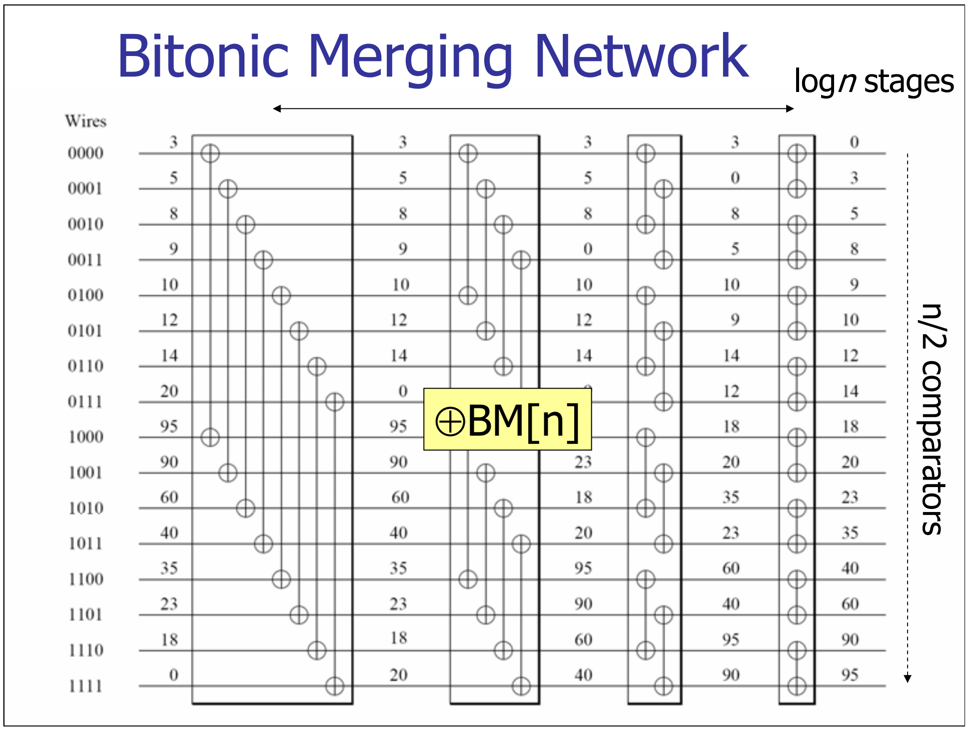
Example: <1,2,4,7,6,0> & <8,9,2,1,0,4> are bitonic sequences.

Bitonic Sort

- Rearrange a bitonic sequence to be sorted.
- Divide & conquer type of algorithm (similar to quicksort) using **bitonic splits**.
  - Sorting a bitonic sequence using bitonic splits = bitonic merge.
  - But we need a bitonic sequence…



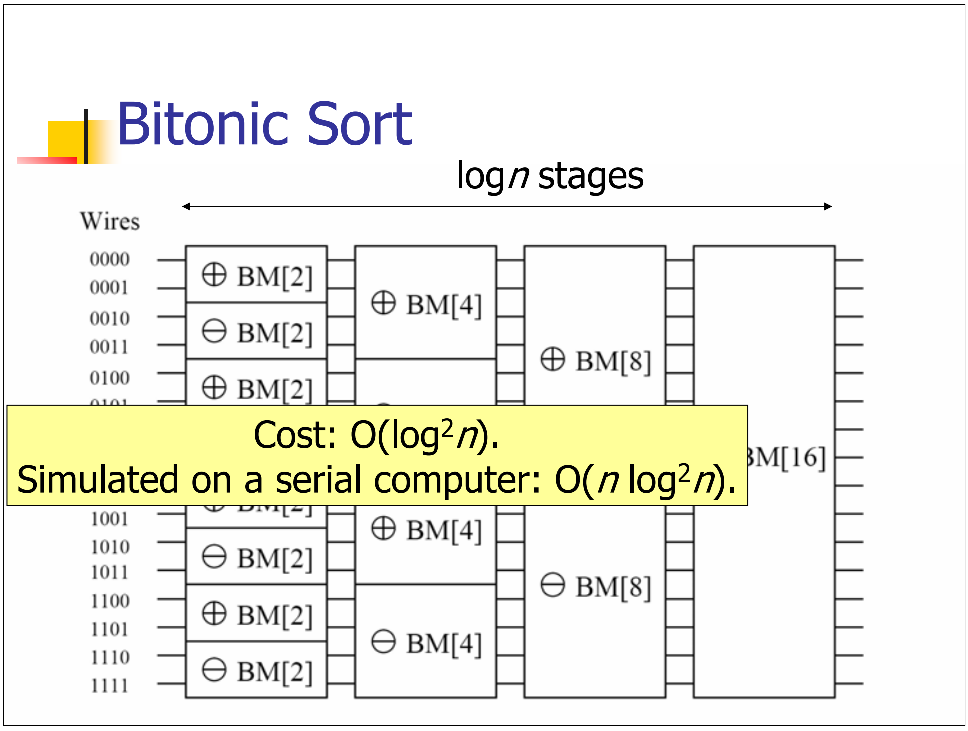
And in fact the procedure works even if the original sequence needs a cyclic shift to look like this particular case.



Cost: Θ(log*n*) obviously.

Bitonic Sort

- Use the bitonic network to merge bitonic sequences of increasing length… starting from 2, etc.
- Bitonic network is a component.



Not cost optimal compared to the optimal serial algorithm.

Mapping to Hypercubes & Mesh – Idea

- Communication intensive, so special care for the mapping.
- How are the input wires paired?
  - Pairs have their labels differing by only one bit $\Rightarrow$ mapping to hypercube straightforward.
  - For a mesh lower connectivity, several solutions but worse than the hypercube $T_P=O(\log^2 n)+O(\sqrt{n})$ for 1 element/process.
  - Block of elements: sort locally ($n/p \log n/p$) & use bitonic merge $\Rightarrow$ cost optimal.

But not efficient & not scalable because the sequential algorithm is suboptimal.

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Hypercube: Neighbors differ with each other by one bit.

Bubble Sort

```
procedure BUBBLE_SORT(n)
begin
   for i := n-1 downto 1 do
      for j := 1 to i do
         compare_exchange(a_j,a_{j+1});
end
```

$\Theta(n^2)$

- Difficult to parallelize as it is because it is inherently sequential.

It is difficult to sort n elements in time log$n$ using n processes (cost optimal w.r.t. the best serial algorithm in n log$n$) but it is easy to parallelize other (less efficient) algorithms.

Odd-Even Transposition Sort

```
procedure ODD-EVEN(n)
begin
    for i := 1 to n do
    begin
        if i is odd then
            for j := 0 to n/2 − 1 do
                compare-exchange(a_{2j+1}, a_{2j+2});
        if i is even then
            for j := 1 to n/2 − 1 do
                compare-exchange(a_{2j}, a_{2j+1});
    end for
end ODD-EVEN
```

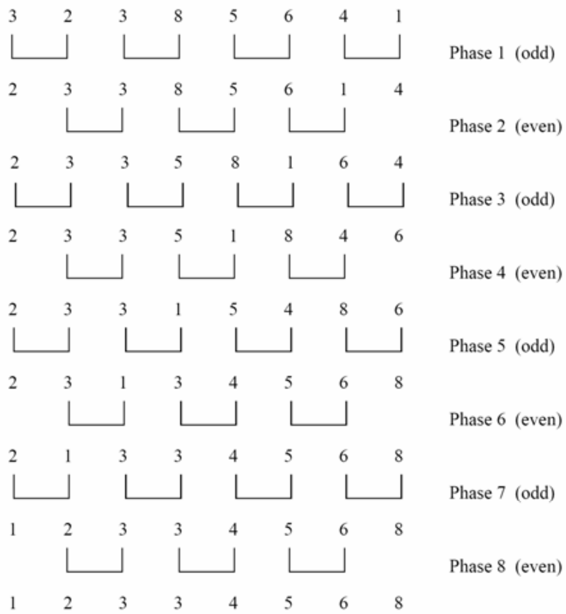
$\Theta(n^2)$

$(a_1,a_2),(a_3,a_4)\ldots$

$(a_2,a_3),(a_4,a_5)\ldots$

Algorithm 9.3 Sequential odd-even transposition sort algorithm.

Unsorted



Sorted

Odd-Even Transposition Sort

- Easy to parallelize!
  - $\Theta(n)$ if 1 process/element.
  - Not cost optimal but use fewer processes, an optimal local sort, and compare-splits:

$$T_P = \Theta\left(\frac{n}{p}\log\frac{n}{p}\right) + \Theta(n) + \Theta(n)$$

local sor ... ation

Cost optimal for $p = O(\log n)$
but not scalable (few processes).

Write speedup & efficiency to find the bound on p but you can also see it with T_P.

Improvement: Shellsort

- 2 phases:
  - Move elements on longer distances.
  - Odd-even transposition but stop when no change.
- Idea: Put quickly elements near their final position to reduce the number of iterations of odd-even transposition.

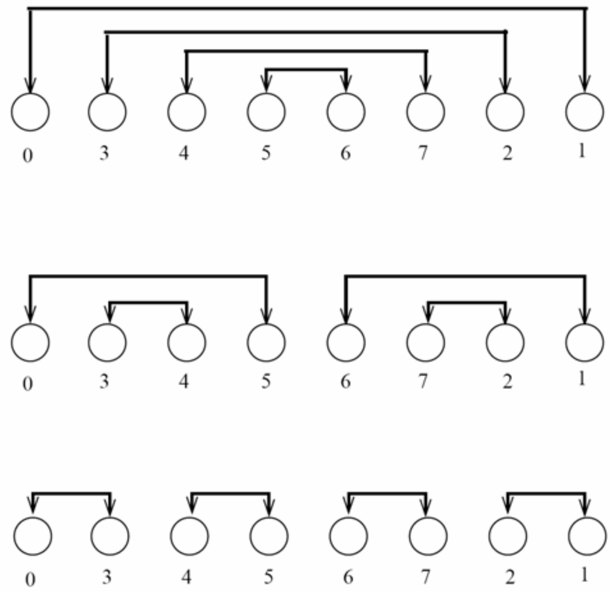
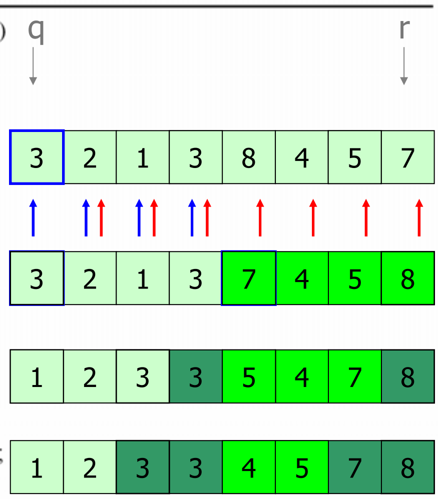


Figure 9.14 An example of the first phase of parallel shellsort on an eight-process array.

Quicksort

- Average complexity: $O(n \log n)$.
  - But very efficient in practice.
  - Average "robust".
  - Low overhead and very simple.
- Divide & conquer algorithm:
  - Partition A[q..r] into A[q..s] ≤ A[s+1..r].
  - Recursively sort sub-arrays.
  - Subtlety: How to partition?

```
procedure QUICKSORT (A, q, r)
begin
    if q < r then
    begin
        x := A[q];
        s := q;
        for i := q + 1 to r do
            if A[i] ≤ x then
            begin
                s := s + 1;
                swap(A[s], A[i]);
            end if
        swap(A[q], A[s]);
        QUICKSORT (A, q, s);
        QUICKSORT (A, s + 1, r);
    end if
end QUICKSORT
```



Algorithm 9.5 The sequential quicksort algorithm.

Hoare partitioning is better. Check in your algorithm course.

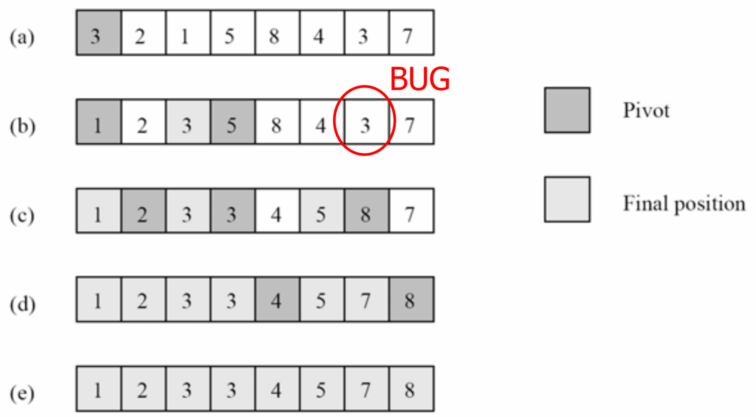


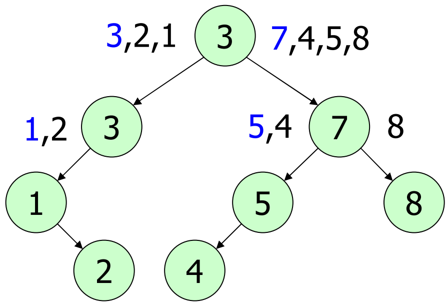
Figure 9.15 Example of the quicksort algorithm sorting a sequence of size $n = 8$.

Parallel Quicksort

- Simple version:
  - Recursive decomposition with one process per recursive call.
  - Not cost optimal: Lower bound = n (initial partitioning).
  - Best we can do: Use $O(\log n)$ processes.
  - Need to parallelize the partitioning step.

Parallel Quicksort for CRCW PRAM

- See execution of quicksort as constructing a binary tree.



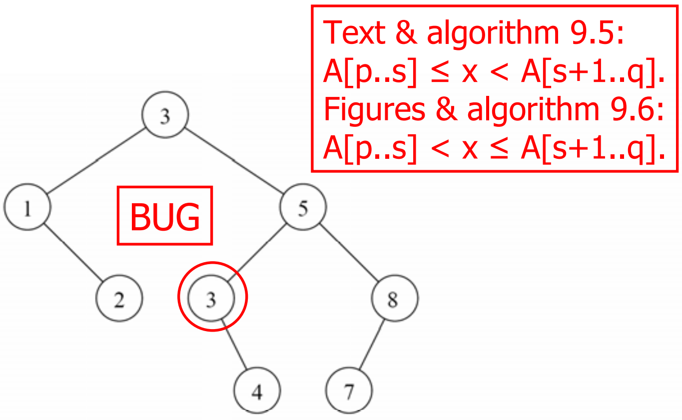


Figure 9.16 A binary tree generated by the execution of the quicksort algorithm. Each level of the tree represents a different array-partitioning iteration. If pivot selection is optimal, then the height of the tree is $\Theta(\log n)$, which is also the number of iterations.

```
procedure BUILD_TREE (A[1 ... n])
begin
    for each process i do
    begin
        root := i;   only one succeeds
        parent_i := root;
        leftchild[i] := rightchild[i] := n + 1;
    end for
    repeat for each process i ≠ root do
    begin
        if (A[i   A[i]≤A[parent_i]
            (A[                        ent_i) then
        begin
            leftchild[parent_i] := i;
            if i = leftchild[parent_i] then exit
            else parent_i := leftchild[parent_i];
        end for
        else
        begin
            rightchild[parent_i] := i;
            if i = rightchild[parent_i] then exit
            else parent_i := rightchild[parent_i];
        end else
    end repeat
end BUILD_TREE
```

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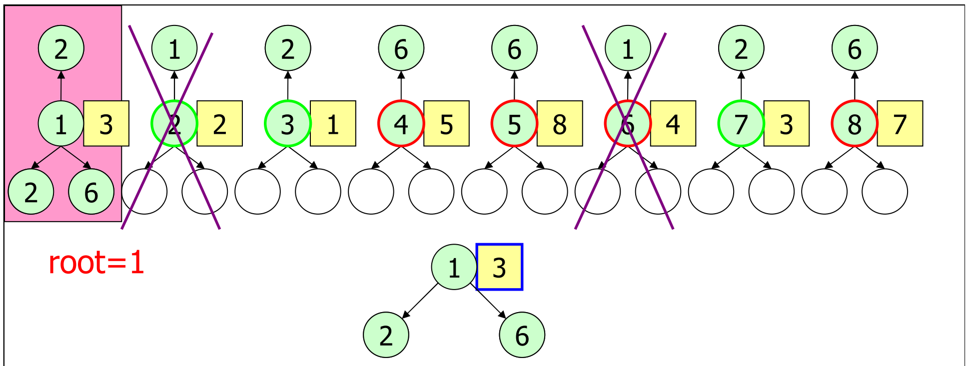
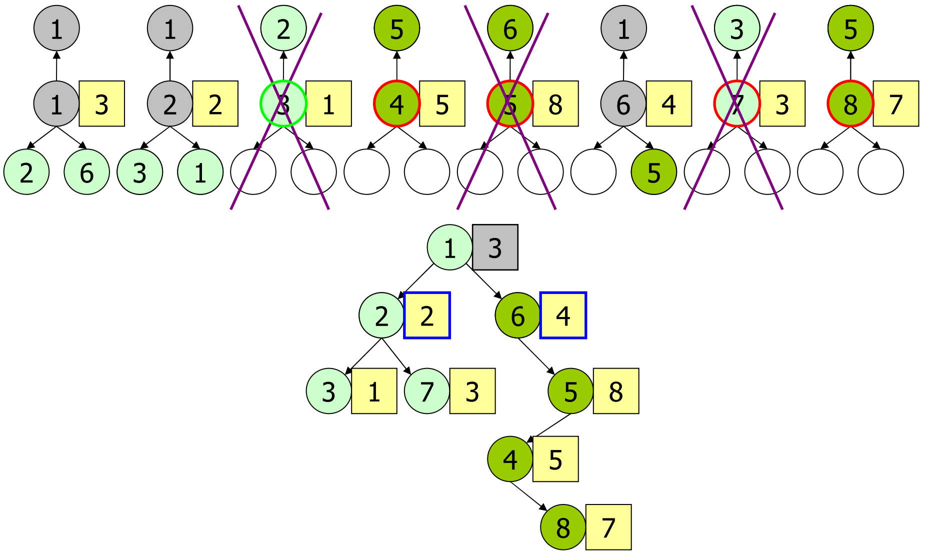
This algorithm does not correspond exactly to the serial version. Time for partitioning: O(1).



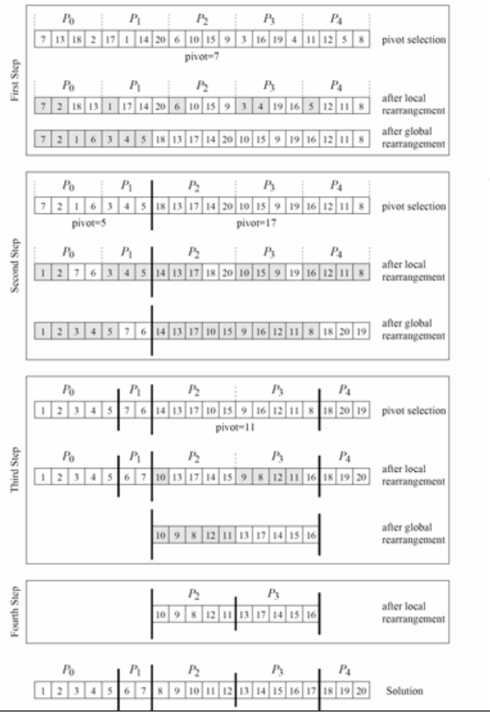
root=1

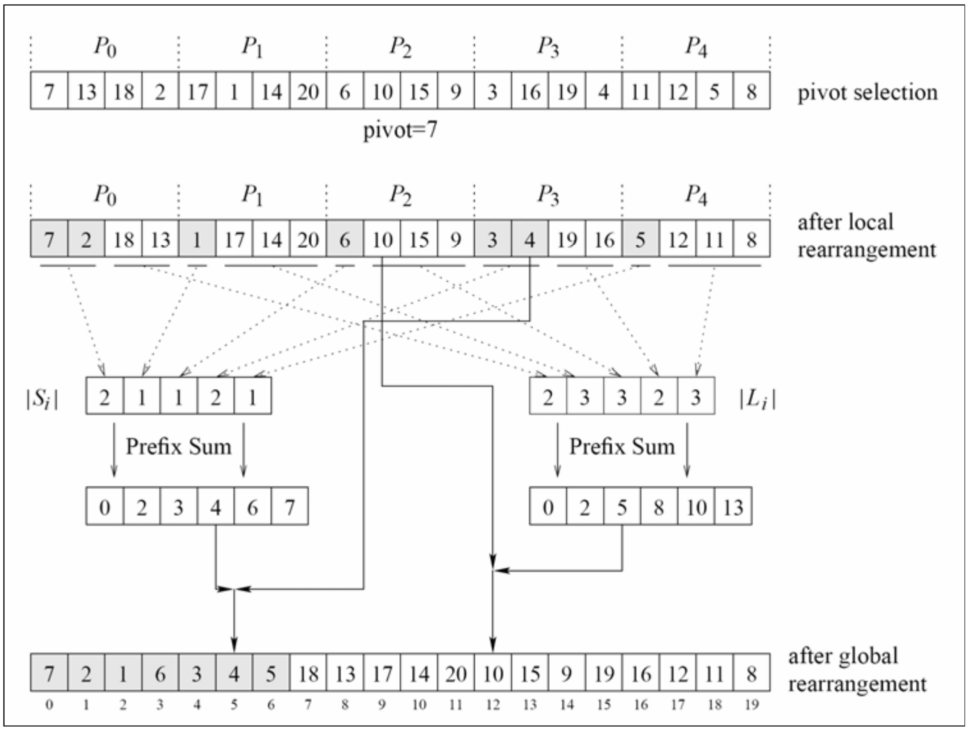


Each step: Θ(1). Average height: Θ(log*n*).
This is cost-optimal – but it is only a model.

Parallel Quicksort – Shared Address (Realistic)

- Same idea but remove contention:
  - Choose the pivot & broadcast it.
  - Each process rearranges its block of elements *locally*.
  - *Global* rearrangement of the blocks.
  - When the blocks reach a certain size, local sort is used.


First Step
pivot selection
pivot=7
after local rearrangement
after global rearrangement
Second Step
pivot=5
pivot=17
Third Step
pivot=11
Fourth Step
Solution



P0 P1 P2 P3 P4
7 13 18 2 17 1 14 20 6 10 15 9 3 16 19 4 11 12 5 8
pivot selection
pivot=7
7 2 18 13 1 17 14 20 6 10 15 9 3 4 19 16 5 12 11 8
after local rearrangement
|Si|
2 1 1 2 1
Prefix Sum
0 2 3 4 6 7
2 3 3 2 3
|Li|
Prefix Sum
0 2 5 8 10 13
7 2 1 6 3 4 5 18 13 17 14 20 10 15 9 19 16 12 11 8
0 1 2 3 4 5 6 7 8 9 10 11 12 13 14 15 16 17 18 19
after global rearrangement

Cost

- Scalability determined by time to broadcast the pivot & compute the prefix-sums.
- Cost optimal.

$$T_P = \overbrace{\Theta\left(\frac{n}{p}\log\frac{n}{p}\right)}^{\text{local sort}} + \overbrace{\Theta\left(\frac{n}{p}\log p\right) + \Theta(\log^2 p)}^{\text{array splits}}. \qquad (4)$$

MPI Formulation of Quicksort

- Arrays must be explicitly distributed.
- Two phases:
  - Local partition smaller/larger than pivot.
  - Determine who will sort the sub-arrays.
    - And send the sub-arrays to the right process.

Final Word

- Pivot selection is very important.
- Affects performance.
- Bad pivot means idle processes.