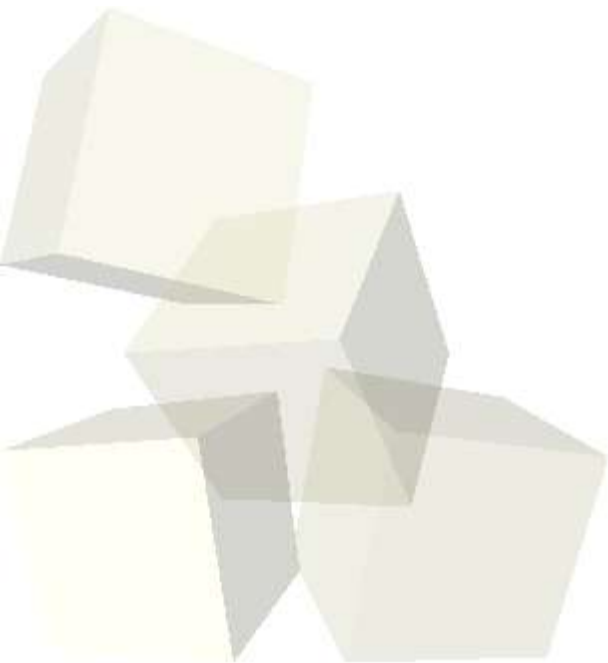
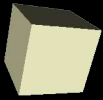


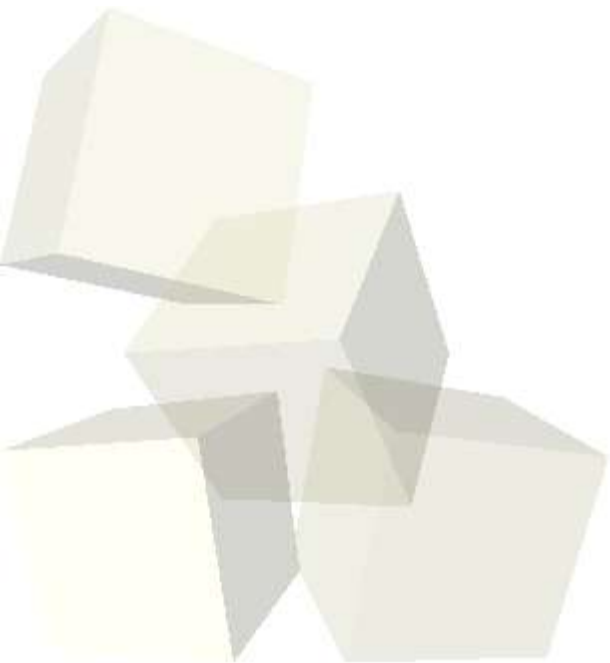
Introduction to Algorithms

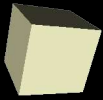
Alexandre David



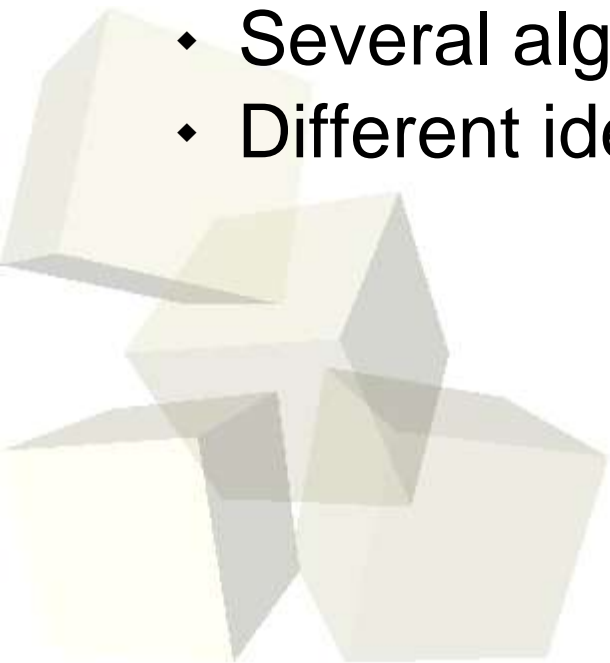


- Notion of algorithms, GCD example.
- Algorithmic problem solving.
- Problem types.
- Sorting example.
- Numerical example.
- Analyzing algorithms.



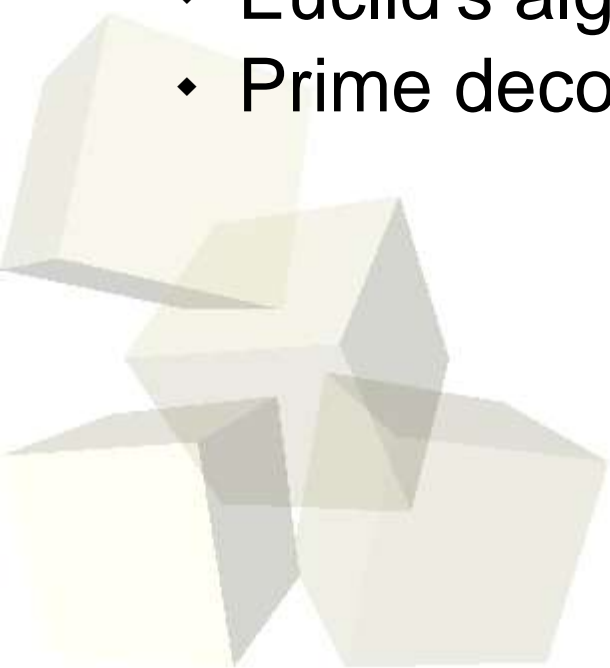


- Why study algorithms?
- What is an algorithm?
- Example: GCD
 - ♦ Known example
 - ♦ Nonambiguity requirement
 - ♦ Define range of inputs
 - ♦ Different representations of the algorithm
 - ♦ Several algorithm for the same problem
 - ♦ Different ideas, different running speeds





- Greatest common divisor of **2 nonnegative, not both zero integers**, denoted $\text{gcd}(m,n)$, defined as the largest integer that **divides both m and n with a remainder of zero**.
- Algorithms:
 - ♦ Consecutive integer checking
 - ♦ Euclid's algorithm
 - ♦ Prime decomposition





Consecutive Integer Checking

- Idea: solution cannot be greater than $\min(m,n)$. Let $t = \min\{m,n\}$. Check t and try again by decreasing t .
- Correctness: greatest? Termination?
- Efficiency: running time?

Step 1: assign $\min\{m,n\}$ to t .

Step 2: divide m by t . If remainder $== 0$ then step 3, otherwise step 4.

Step 3: divide n by t . If remainder $== 0$ return t , otherwise step 4.

Step 4: decrease t by 1, go to step 2.





- Idea: apply repeatedly $\gcd(m, n) = \gcd(n, m \bmod n)$ until $m \bmod n$ is equal to zero (stop when reach $\gcd(m, 0) = m$). Ex: $\gcd(60, 24) = \gcd(24, 12) = \gcd(12, 0) = 12$.

Algorithm *Euclid*(m, n)

// Computes $\gcd(m, n)$ by Euclid's algorithm

// Input: two nonnegative, non both zero integers m and n

// Output: GCD of m and n

while $n \neq 0$ **do**

$r := m \bmod n$

$m := n$

$n := r$

return m



Prime Decomposition

- Idea: decomposition into primes and pick the common factors.

Step 1: find the prime factors of m .

Step 2: find the prime factors of n .

Step 3: identify all the common factors (if p is a common factor occurring p_m and p_n times in m and n , respectively, it should be repeated $\min\{p_m, p_n\}$ times).

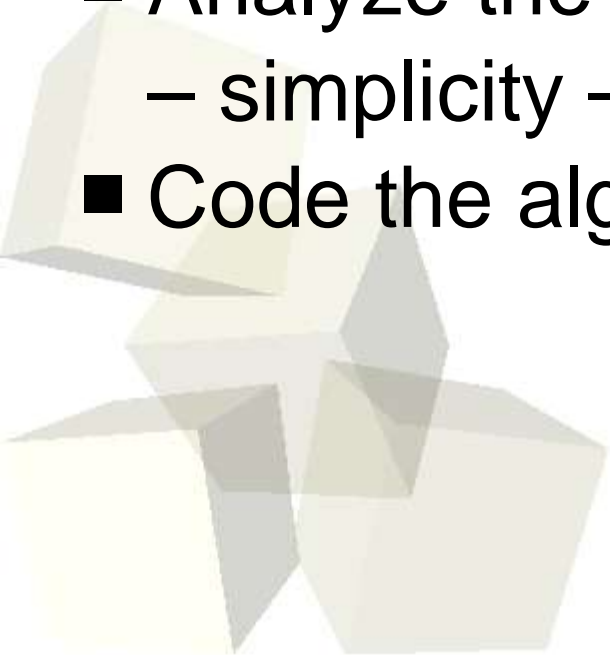
Step 4: Compute the product of all the common factors and return it as the result.

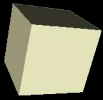
- Problem: Step 1&2 are sub-problems to be solved.



Algorithmic Problem Solving

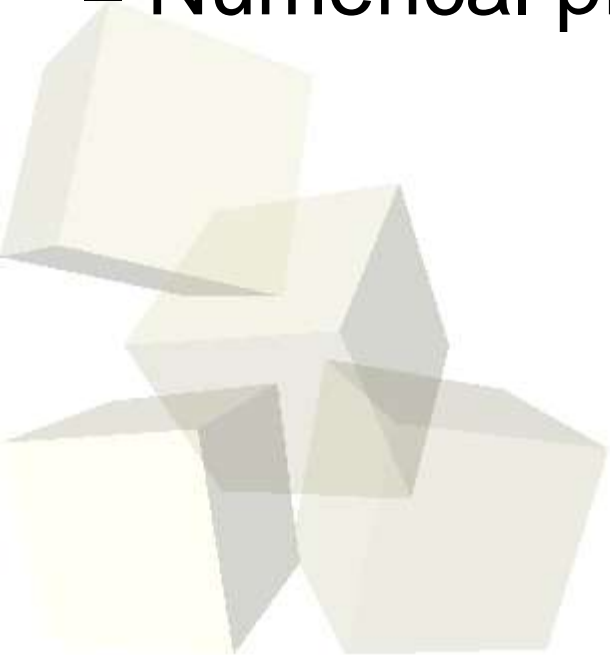
- Understand the problem
- Choose exact/approximate problem solving
- Decide on appropriate data structures
- Apply an algorithm design technique
- Specify the algorithm
- Prove the correctness of the algorithm
- Analyze the algorithm – time and space efficiency
 - simplicity – generality
- Code the algorithm

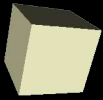




Problem Types

- Sorting
- Searching
- String processing
- Graph problems
- Combinatorial problems
- Geometric problems
- Numerical problems





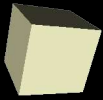
■ The sorting problem:

Input: A sequence of n numbers $\langle a_1, a_2, \dots, a_n \rangle$

Output: A permutation (reordering) $\langle a'_1, a'_2, \dots, a'_n \rangle$ of the input sequence such that $a'_1 \leq a'_2 \leq \dots \leq a'_n$.

■ Algorithms to solve it: insertion sort, merge sort, quicksort. Insertion sort takes $c_1 n^2$ in time, merge sort takes $c_2 n \lg(n)$. Let's sort 10^6 elements.

- ♦ Good insertion code $2n^2$: $2(10^6)^2/10^9 = 2000s$
- ♦ Average merge sort $50n \lg(n)$: $50 * 10^6 \lg(10^6) / 10^7 = 100s$ on another CPU 100x slower.

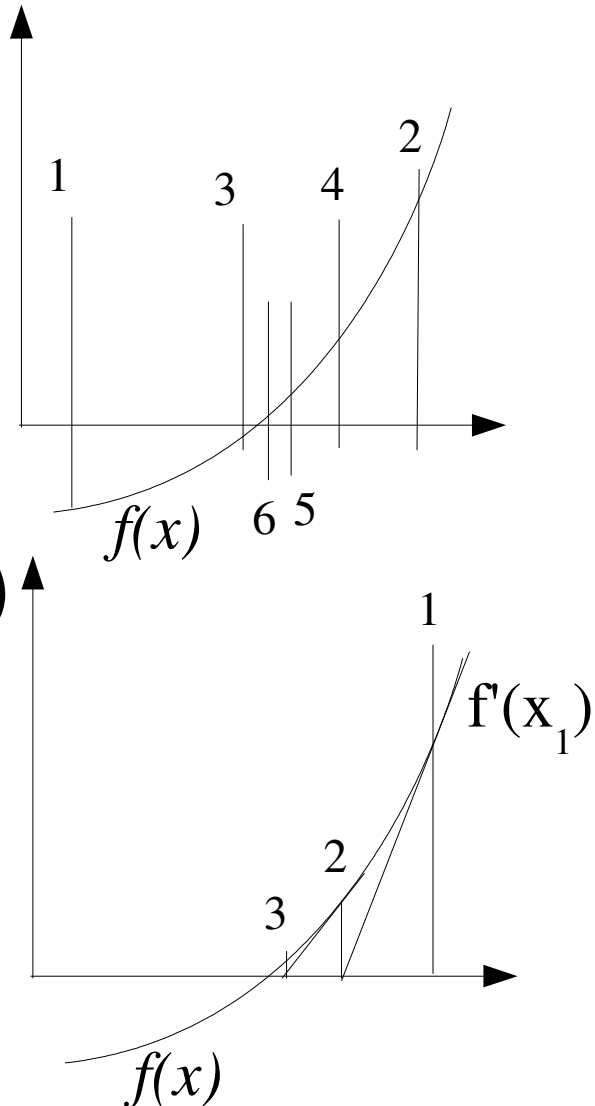


Numerical Example

■ Find x s.t. $f(x)=0$ for a continuous monotonic function.

- Bisection algorithm:
iterate on $[x,y]^0, [x,y]^1 \dots$ s.t. $f(x) < 0$ and $f(y) > 0$ (or opposite), reduce interval by 2 everytime.
- Newton-Raphson algorithm:
use the derivative $x_{i+1} = x_i - f(x_i)/f'(x_i)$
converge much faster.

■ Numerical problems with flat or exponential functions.





- Criteria:
 - ♦ Correctness
 - ♦ Amount of work done
 - ♦ Amount of space used
 - ♦ Simplicity, clarity
 - ♦ Optimality
- Asymptotic behaviour
- Different analysis techniques

